

BusyReplying Bug

- In VENTOS, during the execution of operations, the protocol only considers the messages about the on-going operation.
- Since this logic makes the delivery of unnecessary messages, it needs to be resolved by adding a common message handler.
- Ex) during *Split* operation, only the messages of *Split* micro-operations are considered, and other messages, such as *MERGE_REQUEST* are ignored.

```
else if(vehicleState == state_waitForSplitReply)
{
    if(wsm->getUCommandType() == SPLIT_ACCEPT &&
    ,
else if(vehicleState == state_waitForAck)
{
    if (wsm->getUCommandType() == ACK &&
    {
```

```
else if(vehicleState == state_waitForCHANGEPL)
{
    if (wsm->getUCommandType() == CHANGE_PL &&
    {
else if(vehicleState == state_waitForSplitDone)
{
    if(wsm->getUCommandType() == SPLIT_DONE &&
    {
```

```
else if(vehicleState == state_platoonFollower)
{
    // splitting vehicle receives a SPLIT REQ f
    if (wsm->getUCommandType() == SPLIT_REQ &&
    {
```

MERGE_REQUEST Attempt bug

- When a platoon leader requests *Merge* and the receiver rejects or ignores *MERGE_REQUESTS* more than three times, the sender should no longer send messages.
- However, in VENTOS, since they just change the `mergeReqAttempts` variable to 0, the sender requests Merge again to the same vehicle by the end of the simulation.
- This bug can be solved by storing the ids of vehicles exceeding three times of the *MERGE_REQUESTS*.

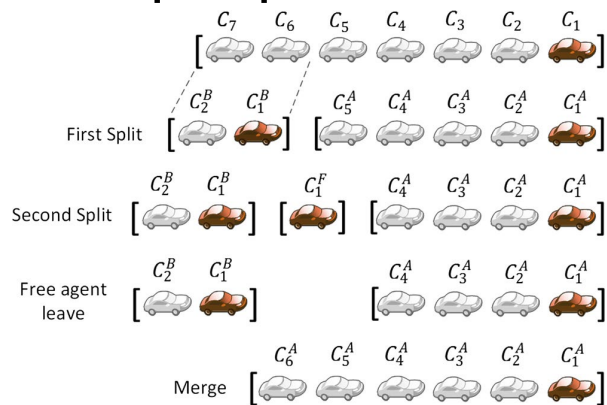
```
if(msg == plnTIMER1) // plnTIMER1: waitForMergeRequest
{
    if(vehicleState == state_waitForMergeReply)
    {
        // leader does not response after three re-attempts!
        if(mergeReqAttempts >= 3)
        {
            mergeReqAttempts = 0;

            setVehicleState(state_platoonLeader);
        }
        else
        {
            setVehicleState(state_sendMergeReq);

            merge_BeaconFSM();
        }
    }
}
```

FollowerLeaveProtocol bug

- In the paper on VENTOS protocol, they described that the *Middle Follower Leave* consists of the sequential operation of *Split*, *Split*, and *Merge*.
- However, the last *Merge* operation is executed by *OptSize* policy and it causes that *Merge* is requested before the end of preceding operations, and thus triggers several unnecessary messages.
- Therefore, it must be implemented that the *Merge* is executed in the proper order according to the specified protocol.



Expected: Split -> Split (Vehicle Leave) -> Merge // FLeave End

Real: Split -> Split (Vehicle Leave) // FLeave End -> Merge

```

// middle follower wants to leave (we need two splits)
else
{
    RemainingSplits = 2;

    // start the first split
    splittingDepth = wsm->getValue().myPltDepth + 1;
    splittingVehicle = plnMembersList[splittingDepth];
    splitCaller = 1; // Notifying split that follower leave is the caller

    setVehicleState(state_sendSplitReq);

    split_DataFSM();
}

```

LeavedVehList bug

- During or after *Leave* operation in VENTOS, we found that a new leader or rear leader sent messages to leaved leader vehicle.
- Since, in VENTOS protocol, when a vehicle leaves a platoon, the vehicle is regarded as an unconnected vehicle, it can not reply to any message.
- Therefore, it is necessary to prevent a vehicle from sending several redundant messages, such as *MERGE_REQUEST*s by *OptSize* policy to leaved vehicles.
- We can solve this bug by storing the ids of vehicles that request *Leave* in *MERGE_REQUEST* attempt code.

LeaveSplitCaller bug (1/2)

- For dealing with various *Leave* operations in VENTOS, their protocol uses splitCaller variable to check the *Leave* cases.
 - Ex) splitCaller – 1: *Follower Leave*, 0: *Leader Leave*, -1: *Split*
- However, in the second *Split* call for *Middle Follower Leave*, the splitCaller needs to be -1 separate from the *End Follower Leave* call.
- For the *Merge* operation to be performed by *Middle Follower Leave*, the leader must accept *MERGE_REQUEST* in a non-busy state after the second split ends.
- But now, since the second *Split* is called with splitCaller=1, the leader keeps busy even after the *Split* ends.

```
else if(RemainingSplits == 1)
{
    // start the second split
    splittingDepth = plnSize - 1;
    splittingVehicle = plnMembersList[splittingDepth];
    splitCaller = 1; // Notifying split that follower leave is the caller

    setVehicleState(state_sendSplitReq);

    split_DataFSM();
}
```

LeaveSplitCaller bug (2/2)

- In current implementation, since the second *Split* is called with `splitCaller=1`, the leader keeps busy even after the *Split* ends.
- To solve this bug, the protocol needs to distinguish the second *Split* call of *Middle Follower Leave* and the first *Split* call of *End Follower Leave*.

```
else if(RemainingSplits == 1)
{
    // start the second split
    splittingDepth = plnSize - 1;
    splittingVehicle = plnMembersList[splittingDepth];
    splitCaller = 1; // Notifying split that follower leave is the caller

    setVehicleState(state_sendSplitReq);

    split_DataFSM();
}
```

ChangeVehStateLastly bug

- In *Split* operation in VENTOS, the last micro-operation to notify that the *Split* is finished is *GAP_CREATED* or *SPLIT_END*.
- However, in the VENTOS protocol, the state of the split vehicle is changed in *CHANGE_PL*.
- The problem is that the busy state of the vehicle is not changed yet.
- Due to this bug, we found quite a few examples of *MERGE_REQUESTS* in the execution of other operations.
- We can solve this bug by synchronizing the state and busy change.

```
if (wsm->getUCommandType() == CHANGE_PL && wsm->
{
    cancelEvent(plnTIMERS5);

    // save my old platoon leader id for future
    oldPlnID = myPlnID;

    // I am a free agent now!
    myPlnID = wsm->getValue().newPltLeader;
    myPlnDepth += wsm->getValue().newPltDepth;
    plnSize = 1;
```